

**PERIODICITY IN THE PERIODIC LAMBDA ALGEBRA**  
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ABSTRACT. The periodic lambda algebra is a co-Koszul complex of the Steenrod algebra whose homology gives the  $E_2$  term for the Adams spectral sequence. Its elements are closely related to periodic homotopy theory, and exhibit periodic properties. In this paper we discuss an algorithm to compute the homology of the periodic lambda algebra, and investigate the algebraic structure for the  $E_2$  page of the  $v_n$  periodic elements.

1. INTRODUCTION

Computing the homotopy groups of spheres is of great interest in algebraic topology. One method of achieving this is by the Adams spectral sequence. This is a spectral sequence whose  $E_2$  term is  $Ext_{A^*}(\mathbb{Z}/p, \mathbb{Z}/p)$ , where  $A^*$  is the Steenrod algebra, that converges to the  $p$ -component of  $\pi_*(S^0)$ .

A well studied object that gives an  $E_1$  term for the Adams spectral sequence is the lambda algebra, which is the co-Koszul complex of the Steenrod algebra by taking the dual to the admissible basis. In [1], Gray gives another algebra  $\Lambda$ , which he terms the periodic lambda algebra, whose homology also gives the  $E_2$  term for the Adams spectral sequence. This is a differential graded algebra generated multiplicatively by elements  $\lambda_i$  and  $v_n$  subject to certain relations.

Compared to the classical lambda algebra, the periodic lambda algebra is smaller and simpler. Of interest to us is that the  $v_n$  generators turn out to correspond to the  $v_n$  self maps in periodic homotopy theory.

It is known that  $H^*(\Lambda_\lambda)$  decomposes into  $v_n$  periodic parts (see §2). This gives us a method to recover  $H^*(\Lambda)$  from  $H^*(\Lambda_\lambda)$ . Our interest here is mainly in determining the algebraic structure of the various periodic parts. We did so by inverting the generator  $v_n$ . An understanding of this can give us information about the homotopy groups of spheres. Since when we delete  $v_1, \dots, v_{n-1}$  and invert  $v_n$ , we obtain the  $E_1$  term of an Adams spectral sequence that converges to  $\pi_*(v_n^{-1}V(n-1))$ . If we then run the Bockstein spectral sequence on the  $E_\infty$  term, we obtain  $\pi_*(S^0)$ .

The project was roughly split into two parts. The first part was writing a computer program to compute the homology of the periodic lambda algebra and

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determine the  $v_n$  (for  $n \leq 3$ ) periodic parts. We did this for the two cases  $p = 3$  and  $p = 5$ . The second part was attempting to understand the algebraic structure of the periodic parts. This was mainly done through comparing the elements with known results and finding patterns between the cases  $p = 3$  and  $p = 5$ . Our main result is determining the generators for the periodic parts, and proving certain elements to be  $v_n$  periodic. Our results also provide some insight into the “telescope conjecture”.

In §2 we give a brief overview of the periodic lambda algebra, which will be the main object of study in our paper. In §3 we discuss our implementation of the leading term algorithm to generate Curtis tables for the periodic lambda algebra, and give a proof of its correctness. In §4, §5 and §7 we present some results we have obtained through investigating the Curtis tables of the respective periodic parts. We also include §6 in which we discuss how our work relates to the telescope conjecture. In §8, we give a theorem for proving certain elements to be  $v_n$  periodic. Finally, in §9 we mention some problems that we wish to continue to work on.

## 2. THE PERIODIC LAMBDA ALGEBRA

The classical lambda algebra is derived from the Steenrod algebra by the taking the Koszul dual of the admissible basis. In [1], Gray derives another co-Koszul complex of the Steenrod algebra by taking the Milnor basis instead, which results in the periodic lambda algebra.

The Milnor basis for the Steenrod algebra (see [3]) is given by the set of generators  $\{Q_k\} \cup \{P^i\}$ , with  $\deg(Q_k) = 2p^k - 1$  and  $\deg(P^n) = 2n(p - 1)$ . These satisfy the relations

(1) If  $a < pb$ , then

$$P^a P^b = \sum (-1)^{a+t} \binom{(p-1)(b-t)-1}{a-pt} P^{a+b-t} P^t$$

(2)  $Q_k Q_l = -Q_l Q_k$

$$(3) P^n Q_k = \begin{cases} Q_k P^n + Q_{k+1} P^{n-p^k} & n > p^k \\ Q_k P^n + Q_{k+1} & n = p^k \\ Q_k P^n & n < p^k \end{cases}$$

According to these relations, we may express all elements in the form

$$Q_{k_1} Q_{k_2} \cdots P^{i_1} P^{i_2} \cdots,$$

where

- (1) all  $Q_k$ 's occur before  $P^i$ 's
- (2) the  $Q_k$ 's are listed in decreasing order
- (3) the portion of the product consisting of  $P^i$ 's are in admissible form.

By the result given in [4], we get a differential graded algebra,  $\Lambda$ , generated by  $\lambda_i$  and  $v_n$  with  $\deg(\lambda_k) = 2k(p - 1)$  and  $\deg(v_n) = 2p^n - 1$  subject to the relations

- (1)  $\lambda_i \lambda_{pi+k} = \sum (-1)^{j+1} \binom{(p-1)(k-j)-1}{j} \lambda_{k+i-j} \lambda_{pi+j}$
- (2)  $v_n v_k = v_k v_n$
- (3)  $\lambda_k v_n = v_n \lambda_k + v_{n-1} \lambda_{k+p^{n-1}}$  if  $n > 0$
- (4)  $\lambda_k v_0 = v_0 \lambda_k$

Homological degree is given by length of monomials. The differentials are given by

- (1)  $d(v_n) = v_{n-1} \lambda_{p^n-1}$  if  $n > 0$
- (2)  $d(v_0) = 0$

$$(3) \quad d(\lambda_k) = \sum (-1)^{j+1} \binom{(p-1)(k-j)-1}{j} \lambda_{k-j} \lambda_j$$

and satisfy the equality

$$d(xy) = d(x)y + (-1)^{\deg(x)} x d(y).$$

Since this is a subcomplex of the cobar complex of the Steenrod algebra, its homology gives  $E_{A^*}(\mathbb{Z}/p, \mathbb{Z}/p)$ .

Now, if we consider the space  $\Lambda/v_0$  arising from the short exact sequence  $0 \rightarrow \Lambda \xrightarrow{v_0} \Lambda \rightarrow \Lambda/v_0 \rightarrow 0$ , we get a Bockstein spectral sequence  $H^*(\Lambda/v_0) \otimes P(v_0) \Rightarrow H^*(\Lambda)$ . We can continue the same process on  $\Lambda/v_0$ , taking the quotient by  $v_1$ , to get  $H^*(\Lambda/(v_1, v_0)) \otimes P(v_0, v_1) \Rightarrow H^*(\Lambda)$ . If we let  $\Lambda_\lambda$  denote the subalgebra of  $\Lambda$  generated by the  $\lambda_i$  terms only, then inductively we can obtain  $H^*(\Lambda_\lambda) \otimes P(v_0, v_1, \dots) \Rightarrow H^*(\Lambda)$ . This shows that  $H^*(\Lambda_\lambda)$  is decomposed into periodic parts.

To find the  $v_n$  torsion free parts of  $\Lambda$ , we first compute  $\Lambda/(v_0, \dots, v_{n-1})$  and then invert  $v_n$ . Computing this and investigating the structure of the periodic parts will be the topic of the remaining of our paper. To simplify notation, we will abbreviate  $v_n^{-1}\Lambda/(v_0, \dots, v_{n-1})$  by  $v_n^{-1}\Lambda$ .

### 3. LEADING TERM ALGORITHM

In [8] Tangora describes an algorithm to compute the homology of the classical lambda algebra. In our project, we modified the algorithm to compute the homology of the periodic lambda algebra. We describe the algorithm and give a proof of correctness.

**3.1. The Algorithm.** We first define an ordering for the terms in  $\Lambda$ . We define all  $v_n$  terms to be greater than  $\lambda_i$ , while  $v_n$  and  $\lambda_i$  are ordered between themselves by topological degree. Monomials in admissible form of the same bidegree are then ordered lexicographically. We say a polynomial is expressed in “proper form” if all monomials are admissible and in decreasing order. Then the order on monomials induces a lexicographic order on polynomials. We will call the greatest monomial in a polynomial the *leading term*.

The algorithm will proceed inductively on total degree and on topological degree for each fixed total degree. At each step, we consider a source and target pair, from a box of bidegree  $(r, s)$  to  $(r-1, s+1)$ . We will say that a proper polynomial  $x$  *tags* a cycle basis  $y$  if it is the smallest polynomial such that  $d(x) = y + (\text{terms} < y)$ .

Suppose a basis at target has already been found. The base case is trivial. For each admissible monomial  $x$  from small to large at source, run the following procedure.

If  $d(x) = 0$ , then add the leading term of  $x$  to the basis at source. Otherwise, let  $m$  be the leading term of  $d(x)$ . If  $m$  has not been tagged, then  $m$  is tagged by the leading term of  $x$ . Otherwise, suppose  $m$  has been tagged by  $y$ . Then replace  $x$  by  $x - y$  and run the above again.

At the end of this process, each basis cycle at target will either remain a basis cycle for the homology group or will be tagged by some element at source in which case both do not belong to the basis. However, the basis elements and tags thus expressed may not necessarily be cycles themselves. They *complete* to *complete cycles* or *full tags*, and are the leading terms of such completions.

The claim is that at each box, the basis is given by the minimal monomial in the homology class it represents.

### 3.2. Proof of correctness.

**Theorem 3.1.** *The algorithm given in §3.1 is finite and correct.*

*Proof.* We first consider finiteness and correctness in the case that we record full cycles and tags. In this case, we are working with complete basis cycles, so we know that at each iteration the leading term of the differential is driven strictly downwards, so the algorithm is finite. It must also return the correct result, as all monomials have either been completed to a full tag or complete cycle. Minimality follows from the fact that the procedure is run in increasing order on the elements at source.

To prove the leading term algorithm correct, we first note two things. First, we may assume that in each box the leading terms of the complete basis cycles are different. For if  $x + y$  and  $x + z$  ( $x > y > z$ ) were both in the basis, we may replace  $x + y$  with  $y - z$  and still obtain the same space. Next, suppose that  $x$  tags  $y$ , but  $d(x)$  does not contain  $y$ . Then the leading term of  $d(x)$  must be larger than  $y$ . This follows from how we arrive at the full tag of a target cycle.

Now, at each iteration in the former case mentioned, when we add in full tags to complete  $x$ , complete basis cycles in  $d(x)$  are reduced in whole. The leading term case of the algorithm can be considered a term-by-term version of this algorithm.

First suppose  $d(x)$  is a single complete basis cycle with leading term  $m$ . Suppose  $m$  is tagged by some element  $t$ . If  $d(t)$  contains  $m$ , then since leading terms of basis cycles are unique, its leading basis cycle must be the complete basis cycle  $d(x)$ . In this case the differential  $d(x - t)$  will be driven downwards and the algorithm will be finite. If  $d(t)$  does not contain  $m$ , then its leading term will be greater than  $m$ . When the procedure is iterated now with  $x - t$  and  $d(x - t)$ , it will behave just as if completing  $t$ . Thus the whole process is just adding in each term of the full tag of the basis cycle one at a time. Since the completion is finite, by induction on the monomials at each source, the process must terminate.

Now for the general case,  $d(x)$  is a combination of complete basis cycles. The algorithm proceeds similarly, adding in terms of full tags. However, in this case, it is possible that the terms of the full tags will overlap each other. This does not have any effect on finiteness, as the worst that can happen is that each term of every full tag in the completion of  $x$  is added. So this reduces to the case where complete cycles and full tags are recorded. Therefore finiteness is proved.

It remains to show that the leading term of the completion of  $x$  is still  $x$ . This is straightforward to check. At any step of the iteration, suppose we determine that  $m$  is already tagged by  $y$ . Then  $y$  must be less than  $x$ , or else the procedure would not have been run on  $y$  yet. This completes our proof. □

An important property the ordering we defined gives rise to is that the initial term of  $d(\mu)\mu_1 \cdots \mu_n$ , where  $\mu$  denotes any generator and  $\mu_1 \cdots \mu_n$  is in admissible form, will be less than  $\mu$ . This is because by the definition of differentials, the leading monomial of the differential of the generators  $\lambda_i$  and  $v_n$  are less than themselves. Since all other terms  $\mu_j$  will be less than or equal to  $\mu$ , the property follows. This property is crucial in the proof of following theorem.

**Theorem 3.2.** *Let  $\mu$  denote a generator. If the monomial  $\mu x$  tags  $\mu y$ , then  $x$  tags  $y$ . Furthermore, if  $x$  tags  $y$ , then  $k\mu x$  tags  $\mu y$  (provided  $\mu x$  and  $\mu y$  are admissible), for some coefficient  $k$ .*

*Proof.* First suppose that  $\mu x$  completes to the tag  $T$  that tags  $\mu y$ . This means that

$$d(T) = \mu y + (\text{terms} < \mu y).$$

On the other hand, we have

$$d(T) = d(\mu)x + \mu d(x) + (\text{smaller terms}).$$

By our previous discussion,  $\mu d(x)$  is the leading term of this polynomial. This implies that

$$d(x) = y + (\text{terms} < y).$$

To prove that  $x$  actually tags  $y$ , we need to show that no term smaller than  $x$  also satisfies this. In order to prove this, we first prove the second part of the theorem.

Suppose  $x$  completes to  $X$ ,  $y$  completes to  $Y$ , and that  $\mu X$  is admissible. Then

$$\mu Y = \mu y + (\text{terms} < \mu Y).$$

$$d(\mu X) = k\mu d(X) + (\text{smaller terms}),$$

where  $k$  is dependent on the degree of  $\mu$ .

Since  $X$  tags  $Y$ , we have as largest term

$$\text{leading term of } k\mu d(X) = k\mu d(x) = \mu y.$$

Now suppose  $\mu S$  is any smaller full tag such that  $d(\mu S) = \mu y + (\text{smaller terms})$ , then by our previous weaker proof of the first statement, we would have that some term less than or equal to  $S$  that tags  $Y$ . But this contradicts our hypothesis that  $X$  tags  $Y$ . Thus the second statement is proved.

Back to the proof of the first statement, suppose that some term  $q < x$  tags  $y$ . Then by the above,  $\mu q$  would tag  $\mu y$ , a contradiction again. This proves the first statement and we are done. □

Due to the parity of degree, in the case of the periodic lambda algebra,  $k = -1$  if  $\mu$  is a  $\lambda$  generator and  $k = 1$  if  $\mu$  is a  $v_n$  generator.

This theorem allows us to suppress terms whose tails have already been tagged, and played an important role in our implementation.

**3.3. Implementation.** We implemented the algorithm using Java. This provided much flexibility in implementation, as Java is an object oriented and relatively low level programming language.

Our first approach was to list out all admissible monomials up to a certain total degree, and run the algorithm on these elements at once. This gave us a Curtis table up to total degree 60. However, after that the RAM space was not enough to store all admissible monomials. This was somewhat expected as the size of the algebra grows exponentially and can be very large.

To circumvent the memory space issue, we re-structured our computer program to generate admissible monomials inductively after the algorithm had been run for the previous total degrees. This allowed us to omit admissible terms whose tails have already been tagged by theorem 3.2. The only resulting issue is that the leading term of a differential may not be in the list of monomials at target. However, if this is the case, it would imply that the leading term is already tagged, so we only need to trace back using theorem 3.2 and find its tag. Under this scheme, the program did not slow down until total degree 130, after which the program naturally became

slow due to the length of each completion, with many terms reaching above 100,000 iterations.

Two days of computing gave us the pretty looking Curtis table for the periodic lambda algebra with highest total degree 146. The terms coincide with the Curtis table for the classical lambda algebra as should. We include the Curtis table up to total degree 50 in the case  $p = 3$ , along with the code of the main method (with certain omissions to increase readability), in the appendix.

We note that computing the homology up to a certain total degree, say  $a$ , does not give all terms for topological degrees up to  $a$ . However, we know there are vanishing lines in the  $E_2$  terms. This is a line of fixed slope, say  $k$ , such that all terms lie below it. Thus, the homology is fully computed at least up to topological degree  $\frac{a}{k+1}$ . Nonetheless, the terms above  $\frac{a}{k+1}$  are still useful in helping to determine generators and structure.

Since we are primarily concerned with the  $v_n$  periodic parts, we also ran a variation of the above program. Rather than first computing the full Curtis table then quotienting out the elements  $v_1, \dots, v_{n-1}$ , it became clear that it would be more efficient to directly quotient the terms during the algorithm. This saved time as the number of elements in each degree decreased and some terms whose completions were very long were omitted. There seemed to be a balance between the number of admissible monomials we omitted and how high our total degree needed to get for us to see periodicity in the table.

We used the computer to find periodic parts in these generated tables. We did so by implementing a method that would take in a confidence factor  $c$ , and search the list for  $v_n$  towers of each cycle. This meant that an element  $x$  would be deemed  $v_n$  periodic if and only if  $v_n^q x$  is a cycle for all  $q$  such that  $v_n^q x$  is in the table, and  $\max\{q\} \geq c$ . For lower degree terms, where  $\max\{q\}$  are relatively large, we can be fairly certain that they are indeed periodic. In fact, many can be proved periodic using theorem 8.1 in §8.

#### 4. $v_1$ PERIODICITY

We computed  $H^*(\Lambda/v_0)$  for  $p = 3$  up to total degree 164 and found the  $v_1$  periodic parts with confidence 5. This gave us a table with highest total degree 129. We also computed the same object up to total degree 303 for  $p = 5$ , which gave us  $v_1$  periodic terms up to degree 247. These gave us enough information to see what happens in lower degrees and to conjecture its general structure for arbitrary  $p$ .

We compared our results to the following theorem, given in [2] by Miller.

**Theorem 4.1.** *The  $E_2$  term for the Adams spectral sequence of  $V(0)$  localizes at  $v_1$  to*

$$P[v_1, v_1^{-1}] \otimes E[h_i : i \geq 1] \otimes P[b_i : i \geq 1]$$

for  $p > 2$ , where the generators are explicit in the cobar complex, with

$$h_i = \{[\bar{\xi}_i]\}, |h_i| = (1, 2(p^i - 1))$$

$$b_i = \left\{ \sum_{j=1}^{p-1} \binom{p}{j} [\bar{\xi}_i^j][\bar{\xi}_i^{p-j}] \right\}, |b_i| = (2, 2p(p^i - 1)).$$

For  $p = 3$ ,  $\lambda_1$  and its Massey product  $\lambda_2 \lambda_1$  are expected generators. The topological degree of  $h_2$  corresponds to the term  $\lambda_4$ . However, since  $\lambda_4$  is not a cycle,

we inverted  $v_1$  to obtain  $v_1^{-1}v_2\lambda_1$  as a candidate for  $h_2$ . Going further down the list,  $v_1^{-3}v_2^3\lambda_2\lambda_1$  at bidegree  $(46, 2)$  gives an element with correct degree to correspond to  $b_2$ . This gave us some clue that the terms  $h_i$  and  $b_i$  correspond to the  $v_2$  towers of  $\lambda_1$  and  $\lambda_2\lambda_1$ . The same pattern was apparent in the  $p = 5$  case, in which case  $b_1$  corresponds to  $\lambda_4\lambda_1$ .

We thus give the following conjecture.

**Conjecture 4.2.** *The generators in theorem 4.1 correspond to the following elements in  $H^*(v_1^{-1}\Lambda)$*

$$h_i \leftrightarrow v_1^{-r}v_2^r\lambda_1$$

$$b_i \leftrightarrow v_1^{-s}v_2^s b_1, \quad b_1 \leftrightarrow \lambda_{p-1}\lambda_1$$

where  $r = \frac{p^{i-1}-1}{p-1}$  and  $s = \frac{p^n-p}{p-1}$ . Thus these terms give the polynomial and exterior generators for the  $v_1$  periodic part in  $\Lambda$ .

It is straightforward to verify that the degrees match, and that  $b_1$  is the Massey product for  $\lambda_1$  as expected.

Each element in the  $v_1$  periodic part can be given a name using these generators. The non-trivial relations for lower degree terms in the case  $p = 3$  are

$$h_1h_2 \leftrightarrow \lambda_2\lambda_3$$

$$h_1h_3 \leftrightarrow v_1^{-3}v_2^3\lambda_2\lambda_3$$

$$h_1h_3b_2 \leftrightarrow v_1^{-1}v_2\lambda_2\lambda_3\lambda_6\lambda_3$$

$$h_2h_3 \leftrightarrow v_1^{-1}v_3\lambda_1\lambda_2\lambda_3.$$

These relations completely define the structure of  $H^*(v_1^{-1}\Lambda)$  up to topological degree 116. Similar relations can be given in the case  $p = 5$ . A diagram for  $p = 3$  in this range is given in figure 1. We omitted the lines indicating multiplication by  $b_2$  to simplify the diagram.

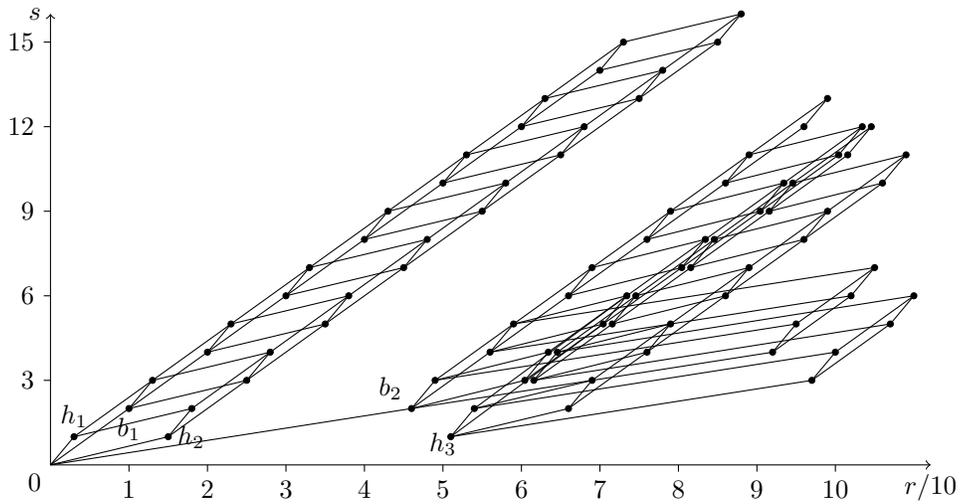


FIGURE 1



The elements  $h_{ij}$  have bidegree  $(i, 2p^j(p^i - 1))$ , and  $b_{ij}$   $(2, 2p^{j+1}(p - 1))$ .

For this case, we computed  $\Lambda/(v_0, v_1)$  up to total degree 149, and found the  $v_2$  periodic parts with confidence 3. This gave us  $v_2$  periodic parts up to topological degree 60. Using this information, we were able to easily identify  $E[h_{10}, h_{11}] \otimes P(b_{10}, b_{11})$ .

**Conjecture 5.2.** *The generators in theorem 5.1 correspond to the generators in  $v_2^{-1}H^*(\Lambda/(v_0, v_1))$  by*

$$\begin{aligned} h_{10} &\leftrightarrow \lambda_1 \\ h_{11} &\leftrightarrow \lambda_3 \\ b_{10} &\leftrightarrow \lambda_2\lambda_1 \\ b_{11} &\leftrightarrow \lambda_6\lambda_3. \end{aligned}$$

Apart from these permanent cycles, we also have the generators  $\lambda_2\lambda_3$ ,  $\lambda_4\lambda_3$  and  $v_2^{-1}v_3\lambda_3$ . However, we are uncertain about the relations these generators satisfy. Although we do know that the relation  $h_{10}h_{11} = 0$  in theorem 5.1 is apparent in the  $E_2$  term.

We have been unable to determine  $\xi$  and  $\zeta_2$  for degree reasons up to our computed range. Also, we have not been able to identify any possible element that can kill  $\lambda_2\lambda_3$  through  $d_2$  or higher differentials. This suggests the possibility that  $\lambda_2\lambda_3$  may possibly be a permanent cycle. These two observations shed some light on the telescope conjecture which we discuss in greater detail in §6.

5.2.  $p > 3$ . Next we consider the  $p = 5$  case. We computed  $H^*(\Lambda/(v_0, v_1))$  up to total degree 304, and found the  $v_2$  periodic part with confidence 3. The element of largest topological degree was at 131.

In [6] Ravenel proves the following theorem.

**Theorem 5.3.** *For  $p > 3$ ,  $H^*(S(2)) \simeq \mathbb{F}_p[v_2, v_2^{-1}]\{1, h_0, h_1, g_0, g_1, h_0g_1\} \otimes E[\zeta]$ , where the bidegrees  $(s, t)$  of the generators are*

$$\begin{aligned} |h_0| &= (1, 2(p - 1)) \\ |h_1| &= (1, -2(p - 1)) \\ |g_0| &= (2, 2(p - 1)) \\ |g_1| &= (2, -2(p - 1)) \\ |\zeta| &= (1, 0). \end{aligned}$$

As with the previous discussions, we have conjectured an identification as follows.

**Conjecture 5.4.** *The elements in theorem 5.3 correspond to elements in the periodic lambda algebra by*

$$\begin{aligned} h_0 &\leftrightarrow \lambda_1 \\ h_1 &\leftrightarrow v_2^{-1}\lambda_5 \\ g_0 &\leftrightarrow v_2^{-1}\lambda_2\lambda_5 \\ g_1 &\leftrightarrow v_2^{-2}\lambda_6\lambda_5 \end{aligned}$$

With this correspondence we have  $h_0g_1 \leftrightarrow \lambda_1\lambda_6\lambda_5 = \lambda_2\lambda_5\lambda_5$ , which completes the “diamond” (see figure 3).

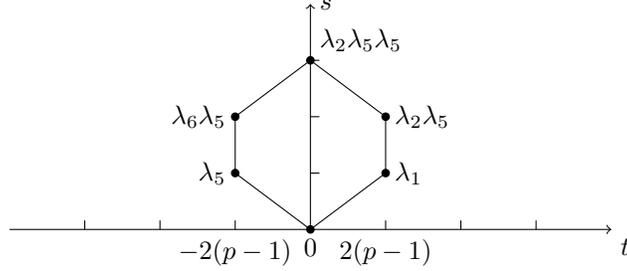


FIGURE 3

In the  $E_2$  page, we also have  $v_3\lambda_5$  and the Massey product  $\lambda_4\lambda_1$  as generators. There are more conjectured generators which we describe in generality in §7.

As with the case  $p = 3$ , we have also been unable to identify the element  $\zeta$  existing in the given algebraic structure. Also, since the terms in our above conjecture gives the permanent cycles, we would expect all other terms not belonging to the diamonds to be killed at some point in the spectral sequence. However, we have been unable to identify these either. We checked the generators  $\lambda_4\lambda_1$  and  $v_3\lambda_5$ , both of which we would expect to be killed. This suggests that  $\lambda_4\lambda_1$  and  $v_3\lambda_5$  may possibly be permanent cycles, in which case the telescope conjecture would also be disproved. We come back to this issue in §6.

## 6. THE TELESCOPE CONJECTURE

As discussed in §5, the telescope conjecture conjectures the equality of  $L_{K(2)}V(1)$  and  $v_2^{-1}V(1)$ . Ravenel conjectured it in [5], then gave a disproof of it in [7]. However, the proof was later deemed incorrect. It remains a conjecture, though most people believe it to be wrong. The main approach in the incorrect proof was to prove the nonexistence of  $\zeta$  in  $v_2^{-1}V(1)$ , which would imply that equality cannot hold, as their  $E_\infty$  terms would not agree. In this aspect, our computations somewhat supports this assertion, despite the proof being incorrect.

For  $p = 3$ , we have been unable to determine  $\zeta_2$  up to total degree 149, and for  $p = 5$  up to 304. Although this gives us evidence of its non-existence up to a certain range, we know the degree of  $\zeta$  may be arbitrarily large. Thus our computations are not sufficient to give a disproof.

Apart from determining the  $\zeta$  terms, we have also run into trouble identifying suitable differentials to kill the non-permanent cycles. This cut-in point seems more promising, as it may be possible to prove the non-existence of differentials, which would also show that the  $E_\infty$  terms do not match.

We consider the generator  $\lambda_2\lambda_3$  in the case  $p = 3$ . First, we know that it is  $v_2$  periodic, by theorem 8.1 which we later prove. On the other hand, we conjecture that differentials should increase the number of  $\lambda$  terms. A verification of differentials that are known shows it is likely that differentials increase  $\lambda$  terms by 1 in general. If this does happen to be the case, then for any term to hit  $\lambda_2\lambda_3$ , it can

only contain one  $\lambda$  term. This would narrow down our search, and could provide a possibility of proving the non-existence by degree means. This is still in work.

### 7. $v_n$ PERIODICITY FOR HIGHER $n$

We first give a general conjecture showing sufficiency in  $v_n$  periodicity.

**Conjecture 7.1.** *If an element is  $v_{n-1}$  periodic and is non-trivial in  $v_n^{-1}\Lambda$ , then it is also  $v_{n+1}$  periodic.*

This holds for  $v_1, v_2$  and  $v_3$  periodicity up to the ranges we have computed.

For higher  $v_n$ , we also observed a neat pattern in the  $p = 3$  case which we also believe to hold for general  $p$ .

From the tables we obtained, for  $v_1$ , we have the generators

$$\lambda_1, v_2\lambda_1, v_2^{13}\lambda_1, \dots$$

$$\lambda_2\lambda_1, v_2^3\lambda_2\lambda_1, v_2^{12}\lambda_2, \dots ;$$

for  $v_2$ , there are the extra generators

$$\lambda_3, v_3\lambda_3, \dots$$

$$\lambda_6\lambda_3, \dots$$

$$\lambda_2\lambda_3 = \langle 1, 1, 3 \rangle, \lambda_4\lambda_3 = \langle 3, 1, 3 \rangle$$

and for  $v_3$ , there is the extra generator  $\lambda_9$  and the conjectured Massey product  $\lambda_{18}\lambda_9$ .

This pattern motivated us to conjecture that in general we have inductive generators for  $v_n^{-1}\Lambda$ .

**Conjecture 7.2.** *The set of generators for  $v_n^{-1}\Lambda$  consists of the  $\lambda$  generators for  $v_{n-1}^{-1}\Lambda$ , in addition to the generators*

$$v_n^{-l}v_{n+1}^l\lambda_{p^{n-1}}$$

$$v_n^{l-1}v_{n+1}^{l-1}\lambda_{2^*p^{n-1}}\lambda_{p^{n-1}}$$

for  $i \geq 1$ , where  $l = \frac{p^{i-1}-1}{2}$ , and the Massey products of the  $\lambda$  terms.

If true, this would be a remarkable result, as it shows much simplicity in the periodic parts of the  $E_2$  term.

### 8. PROOF OF PERIODICITY

In this section, we describe a method to prove that a term is  $v_n$  periodic. This method works for low degree terms which we know to be cycles, and even better if we know that it is  $v_n$  torsion free up to some power of  $v_n$ .

Consider a cycle  $x$  in  $H^*(v_n^{-1}\Lambda)$  with topological degree  $r$ .

We first note that  $v_n^l x$  cannot tag any term for any  $l > 0$ . For if it did, it must tag a term of the form  $v_n^l y$ , as any other term will contain some power of  $v_{n-1}$  which is trivial in  $v_n^{-1}\Lambda$ . But then by theorem 3.2, this would imply that  $x$  tags  $y$ , a contradiction to  $x$  being a cycle.

So if  $x$  is not  $v_n$  periodic, it must be that some element tags an element in the  $v_n$  tower of  $x$ . Let  $l$  be the smallest power such that  $v_n^l x$  is tagged. Then its tag cannot contain any  $v_n$  terms, otherwise it would contradict the minimality of  $l$  by

theorem 3.2. So the smallest possible tag is  $v_{n+1}^l y$  to get the leading generators  $v_n^l$ . In this case, for the correct degrees to be possible, we have the following inequality

$$2l(p^n - 1) + r > 2l(p^{n+1} - 1) + 1.$$

This gives us

$$r - 1 > 2lp^n(p - 1).$$

Suppose we know that  $x$  is  $v_n$  torsion free up to power  $l_0$ . Then the above equality needs to hold for some  $l > l_0$ . If it is the case that  $r - 1 \leq 2(l_0 + 1)p^n(p - 1)$ , then we know that no such term may exist. So it must be that  $x$  is  $v_n$  periodic.

This proves the following theorem.

**Theorem 8.1.** *Suppose  $x$  is of topological degree  $r$  and is such that  $v_n^l x$  is a cycle for all  $l \leq l_0$ . If  $r - 1 \leq 2(l_0 + 1)p^n(p - 1)$ , then  $x$  must be  $v_n$  periodic.*

Recall the generator  $\lambda_2 \lambda_3$  mentioned in §6. Here  $l_0 = 8$  from our Curtis table for  $\Lambda/(v_0, v_1)$  and  $r = 18$ . So  $2(l_0 + 1)p^n(p - 1) = 324$ , while  $r - 1 = 17$ . This proves our claim that  $\lambda_2 \lambda_3$  is  $v_2$  periodic.

## 9. FURTHER WORK

As time was limited in doing this project, and much time was spent on coding the computer program, there are still many problems that we wish to solve. First and foremost is finding proofs (or disproofs) for our conjectures. This will likely require more sophisticated methods, but should be an interesting topic to work on. Ultimately, it would be nice if we could fully determine the algebraic structures for the  $v_n$  periodic parts.

Another promising topic is the telescope conjecture as discussed in §6. If our proposed method does follow through, we would be able to give a disproof of the conjecture for the special cases  $p = 3$  and  $p = 5$ .

We also believe it is worth continue computing the homology of the periodic lambda algebra. Our code could be further optimized and run on faster computers to improve results. This could give us more information to either support or disprove our claims. Another problem that could be considered is determining a general condition for which the leading term algorithm can be applied.

### APPENDIX A. CURTIS TABLE, P=3

We include the Curtis table for  $\Lambda$  when  $p = 3$  up to total degree 50. Generators are represented by numbers, with positive ones  $\lambda$  and negative ones  $v$ . A term such as  $2(-5 \ 1)$  corresponds to  $2v_5 \lambda_1$ . Tags are listed by the format (monomial)/(tag). For instance,  $1(11)/1(2)$  indicates that  $\lambda_1 \lambda_1$  is tagged by  $\lambda_2$ .

3,1 1(1)  
 3,2 1(0 1)/1(-1)  
 6,2 1(1 1)/1(2)  
 7,2 1(-1 1)  
 7,3 1(0 -1 1)/2(-1 -1)  
 10,2 1(2 1)  
 10,3 1(0 2 1)  
 10,4 1(0 0 2 1)/1(-1 -1 1)  
 11,1 1(3)  
 11,2 1(0 3)

11,3  $1(0\ 0\ 3)$   
 11,4  $1(0\ 0\ 0\ 3)/1(-1\ -1\ -1)$   
 13,3  $1(1\ 2\ 1)$   
 13,4  $1(0\ 1\ 2\ 1)/1(-1\ 2\ 1)$   
 14,2  $1(3\ 1)/2(4)$   
 15,2  $1(-1\ 3)/1(-2)$   
 15,4  $1(-1\ -1\ -1\ 1)$   
 15,5  $1(0\ -1\ -1\ -1\ 1)/1(-1\ -1\ -1\ -1)$   
 16,4  $1(1\ 1\ 2\ 1)/1(2\ 2\ 1)$   
 17,4  $1(-1\ 1\ 2\ 1)$   
 17,5  $1(0\ -1\ 1\ 2\ 1)/2(-1\ -1\ 2\ 1)$   
 18,2  $1(4\ 1)/1(5)\ 1(2\ 3)$   
 18,3  $1(0\ 2\ 3)/2(-2\ 1)$   
 19,5  $1(-1\ -1\ -1\ -1\ 1)$   
 19,6  $1(0\ -1\ -1\ -1\ -1\ 1)/2(-1\ -1\ -1\ -1\ -1)$   
 20,4  $1(2\ 1\ 2\ 1)$   
 20,5  $1(0\ 2\ 1\ 2\ 1)$   
 20,6  $1(0\ 0\ 2\ 1\ 2\ 1)/1(-1\ -1\ 1\ 2\ 1)$   
 21,3  $1(3\ 2\ 1)/2(5\ 1)\ 1(1\ 2\ 3)$   
 21,4  $1(0\ 1\ 2\ 3)/1(-1\ 2\ 3)$   
 22,2  $1(3\ 3)/1(6)$   
 22,5  $1(-1\ -1\ -1\ 2\ 1)$   
 22,6  $1(0\ -1\ -1\ -1\ 2\ 1)$   
 22,7  $1(0\ 0\ -1\ -1\ -1\ 2\ 1)/1(-1\ -1\ -1\ -1\ -1\ 1)$   
 23,3  $1(-1\ -2\ 1)$   
 23,4  $1(0\ -1\ -2\ 1)$   
 23,5  $1(0\ 0\ -1\ -2\ 1)\ 1(1\ 2\ 1\ 2\ 1)$   
 23,6  $1(0\ 0\ 0\ -1\ -2\ 1)\ 1(0\ 1\ 2\ 1\ 2\ 1)/1(-1\ 2\ 1\ 2\ 1)$   
 23,7  $1(0\ 0\ 0\ 0\ -1\ -2\ 1)/2(-1\ -1\ -1\ -1\ -1\ -1)$   
 24,4  $1(3\ 1\ 2\ 1)/2(4\ 2\ 1)\ 1(1\ 1\ 2\ 3)/1(2\ 2\ 3)$   
 25,4  $1(-1\ 1\ 2\ 3)/1(-2\ 2\ 1)$   
 25,6  $1(-1\ -1\ -1\ 1\ 2\ 1)$   
 25,7  $1(0\ -1\ -1\ -1\ 1\ 2\ 1)/1(-1\ -1\ -1\ -1\ 2\ 1)$   
 26,2  $1(6\ 1)/2(7)\ 1(4\ 3)$   
 26,3  $1(0\ 4\ 3)/1(-2\ 3)$   
 26,4  $1(-1\ -1\ 2\ 3)$   
 26,5  $1(0\ -1\ -1\ 2\ 3)/1(-1\ -1\ -2\ 1)$   
 26,6  $1(1\ 1\ 2\ 1\ 2\ 1)/1(2\ 2\ 1\ 2\ 1)$   
 27,6  $1(-1\ 1\ 2\ 1\ 2\ 1)$   
 27,7  $1(-1\ -1\ -1\ -1\ -1\ -1\ 1)\ 1(0\ -1\ 1\ 2\ 1\ 2\ 1)/2(-1\ -1\ 2\ 1\ 2\ 1)$   
 27,8  $1(0\ -1\ -1\ -1\ -1\ -1\ -1\ 1)/1(-1\ -1\ -1\ -1\ -1\ -1\ -1)$   
 28,4  $1(4\ 1\ 2\ 1)/1(5\ 2\ 1)\ 1(2\ 1\ 2\ 3)$   
 28,5  $1(0\ 2\ 1\ 2\ 3)/2(-2\ 1\ 2\ 1)$   
 29,3  $1(3\ 2\ 3)/2(5\ 3)\ 1(2\ 3\ 3)$   
 29,4  $1(0\ 2\ 3\ 3)/2(-1\ 4\ 3)$   
 29,7  $1(-1\ -1\ -1\ -1\ 1\ 2\ 1)$   
 29,8  $1(0\ -1\ -1\ -1\ -1\ 1\ 2\ 1)/2(-1\ -1\ -1\ -1\ -1\ 2\ 1)$   
 30,2  $1(7\ 1)/1(8)$

30,5  $1(-1 -1 -1 2 3)$   
 30,6  $1(0 -1 -1 -1 2 3)/2(-1 -1 -1 -2 1) \quad 1(2 1 2 1 2 1)$   
 30,7  $1(0 2 1 2 1 2 1)$   
 30,8  $1(0 0 2 1 2 1 2 1)/1(-1 -1 1 2 1 2 1)$   
 31,3  $1(-1 -2 3)/2(-2 -2)$   
 31,5  $1(3 2 1 2 1)/2(5 1 2 1) \quad 1(1 2 1 2 3)$   
 31,6  $1(0 1 2 1 2 3)/1(-1 2 1 2 3)$   
 31,8  $1(-1 -1 -1 -1 -1 -1 1)$   
 31,9  $1(0 -1 -1 -1 -1 -1 -1 1)/2(-1 -1 -1 -1 -1 -1 -1)$   
 32,4  $1(3 1 2 3)/2(4 2 3) \quad 1(1 2 3 3)/2(2 4 3)$   
 32,7  $1(-1 -1 -1 2 1 2 1)$   
 32,8  $1(0 -1 -1 -1 2 1 2 1)$   
 32,9  $1(0 0 -1 -1 -1 2 1 2 1)/1(-1 -1 -1 -1 -1 1 2 1)$   
 33,3  $1(6 2 1)/2(8 1)$   
 33,4  $1(-1 2 3 3)/1(-2 2 3)$   
 33,5  $1(-1 -2 1 2 1)$   
 33,6  $1(0 -1 -2 1 2 1)$   
 33,7  $1(0 0 -1 -2 1 2 1)/1(-1 -1 -1 -1 2 3) \quad 1(1 2 1 2 1 2 1)$   
 33,8  $1(0 1 2 1 2 1 2 1)/1(-1 2 1 2 1 2 1)$   
 34,2  $1(6 3)$   
 34,3  $1(0 6 3)$   
 34,4  $1(-1 -1 4 3)/2(-2 -2 1) \quad 1(0 0 6 3)$   
 34,5  $1(0 0 0 6 3)$   
 34,6  $1(0 0 0 0 6 3) \quad 1(3 1 2 1 2 1)/2(4 2 1 2 1) \quad 1(1 1 2 1 2 3)/1(2 2 1 2 3)$   
 34,7  $1(0 0 0 0 0 6 3)/1(-1 -1 -1 -1 -2 1)$   
 34,8  $1(-1 -1 -1 -1 -1 -1 2 1)$   
 34,9  $1(0 -1 -1 -1 -1 -1 -1 2 1)$   
 34,10  $1(0 0 -1 -1 -1 -1 -1 -1 2 1)/1(-1 -1 -1 -1 -1 -1 -1 -1 1)$   
 35,1  $1(9)$   
 35,2  $1(0 9)$   
 35,3  $1(0 0 9)$   
 35,4  $1(0 0 0 9)$   
 35,5  $1(0 0 0 0 9)$   
 35,6  $1(-1 1 2 1 2 3)/1(-2 2 1 2 1) \quad 1(0 0 0 0 0 9)$   
 35,7  $1(0 0 0 0 0 0 9)$   
 35,8  $1(-1 -1 -1 1 2 1 2 1) \quad 1(0 0 0 0 0 0 9)$   
 35,9  $1(0 -1 -1 -1 1 2 1 2 1)/1(-1 -1 -1 -1 2 1 2 1) \quad 1(0 0 0 0 0 0 0 9)$   
 35,10  $1(0 0 0 0 0 0 0 0 9)/1(-1 -1 -1 -1 -1 -1 -1 -1 -1)$   
 36,4  $1(6 1 2 1)/2(7 2 1) \quad 1(4 1 2 3)/1(5 2 3) \quad 1(2 2 3 3)$   
 36,5  $1(0 2 2 3 3)/2(-2 1 2 3)$   
 36,6  $1(-1 -1 2 1 2 3)$   
 36,7  $1(0 -1 -1 2 1 2 3)/1(-1 -1 -2 1 2 1)$   
 36,8  $1(1 1 2 1 2 1 2 1)/1(2 2 1 2 1 2 1)$   
 37,3  $1(3 4 3)$   
 37,4  $1(0 3 4 3)/2(-1 6 3)$   
 37,8  $1(-1 1 2 1 2 1 2 1)$   
 37,9  $1(-1 -1 -1 -1 -1 -1 1 2 1) \quad 1(0 -1 1 2 1 2 1 2 1)/2(-1 -1 2 1 2 1 2 1)$   
 37,10  $1(0 -1 -1 -1 -1 -1 -1 1 2 1)/1(-1 -1 -1 -1 -1 -1 -1 2 1)$

38,2  $1(9\ 1)/2(10)\ 1(7\ 3)$   
 38,3  $1(0\ 7\ 3)/2(-1\ 9)$   
 38,6  $1(4\ 1\ 2\ 1\ 2\ 1)/1(5\ 2\ 1\ 2\ 1)\ 1(2\ 1\ 2\ 1\ 2\ 3)$   
 38,7  $1(-1\ -1\ -1\ -1\ -1\ 2\ 3)\ 1(0\ 2\ 1\ 2\ 1\ 2\ 3)/2(-2\ 1\ 2\ 1\ 2\ 1)$   
 38,8  $1(0\ -1\ -1\ -1\ -1\ -1\ 2\ 3)/1(-1\ -1\ -1\ -1\ -1\ -2\ 1)$   
 39,5  $1(3\ 2\ 1\ 2\ 3)/2(5\ 1\ 2\ 3)\ 1(1\ 2\ 2\ 3\ 3)$   
 39,6  $1(0\ 1\ 2\ 2\ 3\ 3)/1(-1\ 2\ 2\ 3\ 3)$   
 39,9  $1(-1\ -1\ -1\ -1\ 1\ 2\ 1\ 2\ 1)$   
 39,10  $1(-1\ -1\ -1\ -1\ -1\ -1\ -1\ -1\ -1\ 1)\ 1(0\ -1\ -1\ -1\ -1\ 1\ 2\ 1\ 2\ 1)/2(-1\ -1\ -1\ -1\ -1\ 2\ 1\ 2\ 1)$   
 39,11  $1(0\ -1\ -1\ -1\ -1\ -1\ -1\ -1\ -1\ 1)/1(-1\ -1\ -1\ -1\ -1\ -1\ -1\ -1\ -1)$   
 40,4  $1(7\ 1\ 2\ 1)/1(8\ 2\ 1)\ 1(3\ 2\ 3\ 3)/1(4\ 4\ 3)$   
 40,7  $1(-1\ -1\ -1\ 2\ 1\ 2\ 3)$   
 40,8  $1(0\ -1\ -1\ -1\ 2\ 1\ 2\ 3)/2(-1\ -1\ -1\ -2\ 1\ 2\ 1)\ 1(2\ 1\ 2\ 1\ 2\ 1\ 2\ 1)$   
 40,9  $1(0\ 2\ 1\ 2\ 1\ 2\ 1\ 2\ 1)$   
 40,10  $1(0\ 0\ 2\ 1\ 2\ 1\ 2\ 1\ 2\ 1)/1(-1\ -1\ 1\ 2\ 1\ 2\ 1\ 2\ 1)$   
 41,3  $1(6\ 2\ 3)/2(8\ 3)$   
 41,4  $1(-1\ 3\ 4\ 3)/1(-2\ 4\ 3)$   
 41,5  $1(-1\ -2\ 1\ 2\ 3)/2(-2\ -2\ 2\ 1)$   
 41,7  $1(3\ 2\ 1\ 2\ 1\ 2\ 1)/2(5\ 1\ 2\ 1\ 2\ 1)\ 1(1\ 2\ 1\ 2\ 1\ 2\ 3)$   
 41,8  $1(0\ 1\ 2\ 1\ 2\ 1\ 2\ 3)/1(-1\ 2\ 1\ 2\ 1\ 2\ 3)$   
 42,2  $1(10\ 1)/1(11)$   
 42,3  $1(-1\ 7\ 3)$   
 42,4  $1(-1\ -1\ 6\ 3)/1(-2\ -2\ 3)\ 1(0\ -1\ 7\ 3)/1(-1\ -1\ 9)$   
 42,6  $1(3\ 1\ 2\ 1\ 2\ 3)/2(4\ 2\ 1\ 2\ 3)\ 1(1\ 1\ 2\ 2\ 3\ 3)/1(2\ 2\ 2\ 3\ 3)$   
 42,8  $1(-1\ -1\ -1\ -1\ -1\ -1\ 2\ 3)$   
 43,5  $1(6\ 2\ 1\ 2\ 1)/2(8\ 1\ 2\ 1)$   
 43,6  $1(-1\ 1\ 2\ 2\ 3\ 3)/1(-2\ 2\ 1\ 2\ 3)$   
 43,7  $1(-1\ -2\ 1\ 2\ 1\ 2\ 1)$   
 44,4  $1(6\ 1\ 2\ 3)/2(7\ 2\ 3)\ 1(4\ 2\ 3\ 3)/2(5\ 4\ 3)\ 1(2\ 3\ 4\ 3)$   
 44,5  $1(0\ 2\ 3\ 4\ 3)/1(-2\ 2\ 3\ 3)$   
 44,6  $1(-1\ -1\ 2\ 2\ 3\ 3)/1(-2\ -2\ 1\ 2\ 1)$   
 45,3  $1(9\ 2\ 1)/2(11\ 1)\ 1(3\ 6\ 3)$   
 45,4  $1(0\ 3\ 6\ 3)$   
 45,5  $1(0\ 0\ 3\ 6\ 3)/2(-1\ -1\ 7\ 3)$   
 46,2  $1(9\ 3)/2(12)$

## APPENDIX B. CODE

The following is the main method of our program to compute the homology of the periodic lambda algebra. We created classes “Monomial” and “Polynomial” that simplified manipulations. We hope the code is self-explanatory enough to read.

---

```

public static void compute() {

    for(int t = 0; t < MAX_ARRAY_SIZE; t++) {

        //find admissible monomials
        for(int i = 0; t - i > 0; i++) {

            int top = i, hom = t - i;

```

```

if(hom == 1) {
    if((top + 1) % (2 * (p - 1)) == 0) {
        ArrayList<Monomial> temp = new ArrayList<Monomial>(1);
        temp.add(new Monomial((top + 1) / (2 * (p - 1))));
        basis[top][hom] = temp;
    }
    else if((top + 2) % 2 == 0){
        double subscript = Math.log((top + 2) / 2) / Math.log(p
        );
        if(subscript - Math.floor(subscript) == 0.0) {
            ArrayList<Monomial> temp = new ArrayList<Monomial>(1)
            ;
            temp.add(new Monomial((int) - subscript));
            basis[top][hom] = temp;
        }
    }
}

else{
    ArrayList<Monomial> temp_box = new ArrayList<Monomial>();
    for(int lambda = 1; 2 * lambda * (p - 1) - 1 < i; lambda
    ++) {
        int lambda_dim = 2 * lambda * (p - 1) - 1;
        for(int item = 0; item < basis[top - lambda_dim][hom -
        1].size(); item++) {
            if(!basis[top - lambda_dim][hom - 1].get(item).
            is_tagged()) {
                if(!basis[top - lambda_dim][hom - 1].get(item).
                tags_something()) {
                    Monomial temp = basis[top - lambda_dim][hom - 1].
                    get(item).clone();
                    temp.add(1, lambda);
                    if(temp.is_admissible())
                        temp_box.add(temp);
                }
            }
            else {
                Monomial this_tags = basis[top - lambda_dim][hom
                - 1].get(item).this_tags().clone();
                this_tags.add(1, lambda);
                Monomial temp = basis[top - lambda_dim][hom - 1].
                get(item).clone();
                temp.add(1, lambda);
                if(temp.is_admissible() && !this_tags.
                is_admissible()) temp_box.add(temp);
            }
        }
    }

    else {
        Monomial tagged_by = basis[top - lambda_dim][hom -
        1].get(item).tagged_by().clone();
        tagged_by.add(1, lambda);
        Monomial temp = basis[top - lambda_dim][hom - 1].
        get(item).clone();
        temp.add(1, lambda);
        if(temp.is_admissible() && !tagged_by.is_admissible
        ()) temp_box.add(temp);
    }
}

```

```

    }
  }
}

for(int v = 0; 2 * (Math.pow(p, v) - 1) < top; v++){
  int v_dim = (int) (2 * (Math.pow(p, v) - 1));
  for(int item = 0; item < basis[top - v_dim][hom - 1].
    size(); item++) {
    if(!basis[top - v_dim][hom - 1].get(item).is_tagged()
      ) {
      if(!basis[top - v_dim][hom - 1].get(item).
        tags_something()){
        Monomial temp = basis[top - v_dim][hom - 1].get(
          item);
        temp.add(1, -v);
        if(temp.is_admissible()) temp_box.add(temp);
      }
      else {
        Monomial this_tags = basis[top - v_dim][hom - 1].
          get(item).this_tags();
        this_tags.add(1, -v);
        Monomial temp = basis[top - v_dim][hom - 1].get(
          item).clone();
        temp.add(1, -v);
        if(temp.is_admissible() && !this_tags.
          is_admissible()) temp_box.add(temp);
      }
    }
  }

  else {
    Monomial tagged_by = basis[top - v_dim][hom - 1].
      get(item).tagged_by();
    tagged_by.add(1, -v);
    Monomial temp = basis[top - v_dim][hom - 1].get(
      item).clone();
    temp.add(1, -v);
    if(temp.is_admissible() && !tagged_by.is_admissible
      ()) temp_box.add(temp);
  }
}
}
temp_box = sort(temp_box);
basis[top][hom] = temp_box;
}

//curtis
if(t != 4){
  for(int i = 4; t - i > 0; i++) {

    int top = i, hom = t - i;

    for(int item = basis[top][hom].size() - 1; item > -1;
      item--) {

      Polynomial differential = Calculator.
        compute_differential(basis[top][hom].get(item));
    }
  }
}

```

