

# NEW TOPIC: HOMOLOGY

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WE HAVE DEFINED

$\pi_1(X, x_0)$  = GROUP COMPUTABLE  
IN PRACTICE

$\pi_k(X, x_0)$  = ABELIAN GROUP FOR  $k > 1$   
HARD TO COMPUTE

DEFINED IN TERMS OF MAPS

$(\mathbb{I}^k, \partial \mathbb{I}^k) \rightarrow (X, x_0)$  WITH

BINARY OPERATION SIMILAR

TO  $k=1$  CASE

WILL DEFINE

ABSOLUTE  $H_k(X)$  =  $k$ TH HOMOLOGY GROUP OF  $X$   
= ABELIAN GROUP FOR  $k \geq 0$

HARDER TO DEFINE BUT  
EASIER TO COMPUTE WITH.

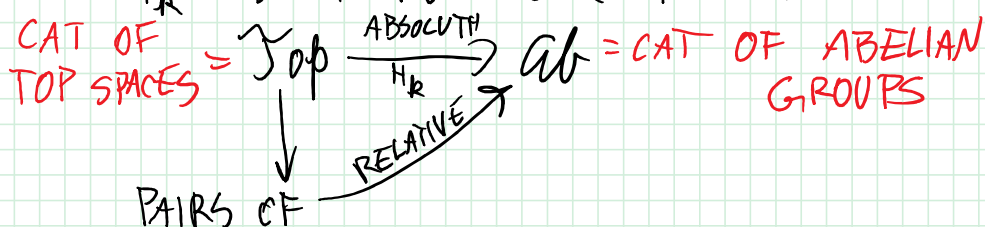
IT HAS A RELATIVE VERSION

$H_k(X, A)$  FOR  $A \subseteq X$

WITH  $H_k(X) = H_k(X, \emptyset)$

IT WILL BE SHOWN THE  
EILENBERG - STEENROD AXIOMS

$H_k$  IS A FUNCTOR FROM



## 1) HOMOTOPY AXIOM

SUPPOSE  $X \xrightarrow{f, g} Y$  THAT ARE HOMOTOPIC

THEN  $H_k(X) \xrightarrow{f_*, g_*} H_k(Y)$  ARE EQUAL

SIMILARLY  $(X, A) \xrightarrow{f_*, g_*} (Y, B)$   
FOR

## 2) EXACTNESS AXIOM

GIVEN  $A \xrightarrow{i} X$ , WE GET

$$\begin{array}{ccccccc} H_k(X, A) & \xrightarrow{j_{k+1}} & H_k(A) & \xrightarrow{i_*} & H_k(X) & \xrightarrow{j_*} & H_k(X, A) & \xrightarrow{\partial_k} & H_{k-1}(A) \rightarrow \dots \\ & & & & \parallel & & & \uparrow & \\ & & & & H_k(X, \emptyset) & & & & \\ & & & & (X, \emptyset) \xrightarrow{t} (X, A) & & & & \end{array}$$

TO BE DEFINED LATER

THIS IS A LONG EXACT SEQUENCE

## 3) EXCISION AXIOM

LET  $B \subset A \subset X$  WITH

CLOSURE(B)  $\subseteq$  INTERIOR OF A

$$H_k(X-B, A-B) \xrightarrow{\cong} H_k(X, A)$$

## 4) DIMENSION AXIOM

$$H_k(\text{point}) = \begin{cases} \mathbb{Z} & \text{FOR } k=0 \\ 0 & \text{FOR } k>0 \end{cases}$$

TWO WAYS TO PROCEED:

① PROVE SOME THINGS  
ASSUMING THESE AXIOMS

② CONSTRUCT A FUNCTOR  
THAT SATISFIES THEM.

WILL PROCEED WITH ② FOR NOW.

WE NEED SOME ALGEBRAIC DEFINITIONS

DEF A CHAIN COMPLEX  $C$   
(OF ABELIAN GROUPS) IS  
A DIAGRAM

$$C_0 \xleftarrow{d_1} C_1 \xleftarrow{d_2} C_2 \xleftarrow{d_3} C_3 \xleftarrow{\dots} \dots$$

EACH  $C_i$  IS AN ABELIAN GRP  
"  $d_i$  IS HOMOMORPHISM

WITH  $d_i d_{i+1} = 0$  FOR  $i \geq 1$ .

(WEAKER THAN EXACTNESS)

$x \in C_n$  IS AN  $n$ -CHAIN

IF  $d_n(x) = 0$ ,  $x$  IS AN  $n$ -CYCLE

IF  $x \in \text{Im } d_{n+1}$ ,  $x$  IS AN  $n$ -BOUNDARY

$d_n = n$ TH BOUNDARY OPERATOR

ITS  $n$ TH HOMOLOGY GROUP IS

$$C_n \supset \text{ker}(d_n) \supset \text{im}(d_{n+1})$$

$$\text{ker}(d_n) / \text{im}(d_{n+1}) =: H_n(C)$$

A CHAIN MAP  $C \xrightarrow{f} C'$

IS A COMMUTATIVE DIAGRAM

$$\begin{array}{ccccccc} C_0 & \xleftarrow{d_1} & C_1 & \xleftarrow{d_2} & C_2 & \xleftarrow{d_3} & C_3 \xleftarrow{\dots} \dots \\ b_0 \downarrow & & b_1 \downarrow & & b_2 \downarrow & & b_3 \downarrow \\ C'_0 & \xleftarrow{d'_1} & C'_1 & \xleftarrow{d'_2} & C'_2 & \xleftarrow{d'_3} & C'_3 \xleftarrow{\dots} \dots \end{array}$$

$$f_n d_{n+1} = d'_{n+1} f_{n+1} \quad \text{FOR } n \geq 0$$

EXERCISE: SHOW THAT  $f$

ABOVE INDUCES HOMOMORPHISMS

$$H_n(C) \xrightarrow{f_*} H_n(C')$$

FOR  $n \geq 0$ .

THEOREM SUPPOSE WE HAVE  
A SHORT EXACT SEQUENCE  
OF CHAIN COMPLEXES

$$0 \rightarrow C' \xrightarrow{f} C \xrightarrow{g} C'' \rightarrow 0$$

i.e.

$$0 \rightarrow C'_n \xrightarrow{f_n} C_n \xrightarrow{g_n} C''_n \rightarrow 0$$

IS EXACT,  $\text{ker } f_n = 0$ ,

$\text{im } f_n = \text{ker } g_n$  AND  $g_n$  IS ONTO  $\forall n \geq 0$

THEN THERE IS A LONG EXACT  
SEQUENCE

CONNECTING HOMOMORPHISM

$$\dots \rightarrow H_n(C') \xrightarrow{f_*} H_n(C) \xrightarrow{g_*} H_n(C'') \xrightarrow{\partial_n} H_{n-1}(C') \rightarrow \dots$$

PROOF

$$\begin{array}{ccccccc}
 0 & \rightarrow & C'_n & \xrightarrow{f_n} & C_n & \xrightarrow{g_n} & C''_n \rightarrow 0 \\
 & & \downarrow d'_n & & \downarrow d_n & & \downarrow d''_n \\
 0 & \rightarrow & C'_{n-1} & \xrightarrow{f_{n-1}} & C_{n-1} & \xrightarrow{g_{n-1}} & C''_{n-1} \rightarrow 0 \\
 & & \downarrow d'_{n-1} & & \downarrow d_{n-1} & & \downarrow d''_{n-1} \\
 0 & \rightarrow & C'_{n-2} & \xrightarrow{f_{n-2}} & C_{n-2} & \rightarrow & 0
 \end{array}$$

$\alpha \in H_n(C'') = \text{ker}(d''_n) / \text{im}(d''_{n+1})$   
 $x \in \text{ker}(d''_n) \subset C''_n$  OF  $\alpha$

HOW TO DEFINE  $\partial_n$ :

$$\alpha \in H_n(C'') = \text{ker}(d''_n) / \text{im}(d''_{n+1})$$

CHOOSE A REPRESENTATIVE

$$x \in \text{ker}(d''_n) \subset C''_n \quad \text{OF } \alpha$$

$x \in \ker(d_n'') \subset C_n''$  OF  $\alpha$   
CHOOSE  $y \in C_n$  WITH  $g_n(y) = x$

NOTE  $g_{n-1} d_n(y) = d_n'' g_n(y) = d_n''(x) = 0$

SO  $d_n(y) \in \ker g_{n-1} = \text{im } f_{n-1}$

AND  $d_n(y) = f_{n-1}(z)$  FOR A UNIQUE  $z$ .

NOTE  $f_{n-2} d_{n-1}'(z) = d_{n-1} f_{n-1}(z) = d_{n-1}(d_n(y)) = 0 \in C_{n-2}$

THIS MEANS  $d_{n-1}'(z) = 0$

$z$  IS AN  $(n-1)$ -CYCLE REPRESENTING  
SOME  $\gamma \in H_{n-1}(C')$

WE NEED TO SHOW  $\gamma$  IS  
INDEPENDENT OF THE TWO  
CHOICES MADE.

A DIFFERENT REPRESENTATIVE  $\tilde{x}$  OF  
 $\alpha$  WOULD HAVE THE FORM

$\tilde{x} + d_{n+1}''(w)$  FOR  $w \in C_{n+1}''$   
 $\parallel$   
 $g_{n+1}(u)$  FOR  $u \in C_{n+1}$

COULD REPLACE  $y$  BY  $y + d_{n+1}(u)$ ,  
THEN  $d_n(y + d_{n+1}(u))$   
 $= d_n(y) + d_n d_{n+1}(u) = d_n(y)$

A DIFFERENT CHOICE OF  $y$   
HAS THE FORM  $y + f_n(v)$   $v \in C'_n$

$$d_n(y + f_n(v)) = d_n(y) + d_n f_n(v) \\ = d_n(y) + f_{n-1} d'_n(v)$$

THIS IS IN  $\ker g_{n-1}$ .

WE WOULD REPLACE  $z$  BY  
 $z + d_{n-1}(v)$

IT DIFFERS FROM  $z$  BY A  
BOUNDARY, SO IT REPRESENTS  
THE SAME  $\gamma \in H_{n-1}(C')$

$\gamma$  IS WELL DEFINED

DIAGRAM CHASING

# MINIMAL CHASING